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Transmitted herewith for filing is the patent application of:

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For: UNIFICATION OF A PROGRAMMING LANGUAGE AND A DEFINITION
LANGUAGE

Enclosed are:

- ☒ 32 pages of specification, 9 pages of claims, an abstract and a Combined Declaration and Power of Attorney.
- ☒ 15 sheet(s) of formal drawings.
- ☒ An assignment of the invention to: Microsoft Corporation and a Recordation Cover Sheet.
- ☒ Information Disclosure Statement.
- ☒ Form PTO-1449 and copies of documents listed thereon.

FILING FEE					
For	Claims Filed	Number Free	Number Extra	Rate	Basic Fee
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TECHNICAL FIELD

The present invention pertains to unifying a programming language with a
definition language. In a compiler environment, a compiler recognizes definition
15 language constructs embedded in programming language source code, creates a
parse tree that unifies representation of the definition language and programming
language, and generates output code based upon the semantics of the definition
language and programming language.

20

BACKGROUND OF THE INVENTION

A complex computer program can contain millions of lines of code. To
manage this complexity, many programmers have adopted object-oriented
programming techniques. Object-oriented programming splits a computer program

into modules called objects. Each object performs specific functions and interacts with other objects in a pre-defined way. Conventionally, programmers use a definition language to create a specification for an object. This definition language specification defines how the object interacts with other objects. With reference
5 to this definition language specification, programmers then use a programming language to actually write code for the object.

Figure 1 shows several principles of object-oriented programming with reference to an object 100 that interacts in pre-defined ways with a client 140 (which can also be an object).

10 The object 100 encapsulates data 110, which indicates the current state of the object 100. The object 100 exposes member functions 120, 122, and 124 that provide access to data 110 or provide some other defined function. To access the data 110, the client 140 of the object 100 goes through a member function. In Figure 1, the member functions 120, 122, 124 perform related
15 services for the client 140, and are grouped into an interface 126. In some object-oriented programming models, an object can expose multiple interfaces. For example, Figure 1 also includes an interface 130 (shown without member functions for simplicity).

A particular member function might take arguments as input and/or
20 produce arguments as output. An argument can be, for example, a number, a character, a string of characters, a user-defined data structure, or a pointer. To interact with an object exposing a member function, a client gives arguments of the correct type and in the correct order. The object processes the data for the

input arguments according to the member function, and then returns output arguments of the correct type in the correct order. A definition language specification for the object 100 defines how the object 100 and client 140 interact, for example, what arguments to exchange, what order to exchange them, etc. The client 140 and the object 100 interact per the definition language specification, after a programmer has written programming language code to implement the client 140 and the object 100 according to the definition language specification.

Definition languages have evolved to precisely specify information about interfaces, the member functions of interfaces, and the objects that expose interfaces. A definition language for interfaces is commonly termed an interface definition language ["IDL"]. Historically, interface definition languages have sometimes been termed object definition languages or object description languages.

Figures 2a and 2b show a source code listing of an IDL file 200 for an object CTest that exposes the interfaces ITest and ITest2. The IDL portion 210 indicates that ITest has methods Grade and Score, and indicates data types for the input and output arguments of Grade and Score. The IDL portion 230 defines ITest2, which includes the methods Display and Hours as well as methods to retrieve and set a value for the property StudentID. The IDL portion 250 describes an object CTest that exposes the ITest and ITest2. This IDL file shows the complexity and intricacy of specifying interfaces with IDL. For more information about the IDL used in IDL file 200, see Kraig Brockschmidt, Inside

OLE, second edition, "Chapter 3, Type Information," Microsoft Press (1995) or Al Major, COM IDL & Interface Design, Wrox Press (1999).

Conventionally, an IDL defines interfaces and objects, while a programming language is used to implement the interfaces for an object according to the IDL specification for the object. IDLs have developed separately from programming languages, follow different rules, and use different tools.

Figure 3 illustrates separate treatment of IDL and programming language code in a "build process" for creating a binary file 380. In Figure 3, an IDL file 300 includes an IDL specification for an object and the interfaces of the object.

10 An IDL compiler 310 converts the IDL file 300 to output code such as a type library file 320 describing information types for the interfaces. The IDL compiler 310 also produces a header file 330 for consumption by the C++ compiler 350.

A C++ source code file 340 includes statements in the programming language C++. In this example, the C++ file 340 is for implementing a class

15 that exposes the interfaces defined in the IDL file 300. The C++ compiler 350 converts the C++ file 340 to the output code 360. The C++ compiler uses the header file 330 when converting the C++ source code to the output code 360.

The type library is packaged in a RC file 324, from which a resource compiler 326 produces a RES file 328 that is passed to the linker 370. The linker 370 packages

20 the RES file 328 together with the output code 360 and any other appropriate resources into a binary file 380 such as an executable or dynamic link library.

Separate treatment of IDL and programming language code has several disadvantages.

First, when an IDL compiler converts IDL to output code, the IDL compiler strips most IDL information. This stripped IDL information is unavailable to a programming language compiler subsequently receiving the output code, which prevents the programming language compiler from fully utilizing combined
5 knowledge of IDL and programming language code.

Second, separately handling IDL and programming language code complicates programming by requiring separate files and more code from a programmer, and by using different rules and tools.

Third, inconsistencies between IDL files and programming language files
10 can create bugs when a project is being built. Worse yet, run time inconsistencies between programming language code and corresponding IDL can cause a program to inexplicably crash.

Microsoft Corporation's Visual C++ 5.0 and 6.0 provide development tools that somewhat simplify separate treatment of IDL. In Visual C++ 6.0, a
15 "ClassWizard" development tool allows a programmer to add features to an object by manipulating a high level representation of the object. When a programmer adds a feature, IDL is automatically generated and added to a separate IDL file. For more information about development tools that facilitate separate treatment of IDL code and C++ code, see Sphar et al., Learn Microsoft Visual C++ 6.0 Now,
20 Microsoft Press (1998) or Horton, Beginning Visual C++ 6, Wrox Press Ltd. (1998).

These development tools still involve an IDL compiler that converts IDL to output code, stripping most IDL information in the process. A programming

language compiler that later handles the output code cannot utilize this lost information. Moreover, although these development tools simplify the separate handling of IDL and programming language code in many respects, they still expose the programmer to separate programming language and IDL files.

5 Other programming tools recognize selected, limited forms of IDL information in programming language source code. These programming tools typically process the IDL information, pass it through to an output file as IDL metadata, and then discard the IDL information. These tools do not derive any semantic meaning from the IDL information.

10 For example, Microsoft Corporation's Visual J++ allows placement of IDL attributes of selected, limited types in comments in Java source code. A pre-processor strips away the comments, but preserves the IDL attributes. Later, a compiler passes the IDL information through to an output file as IDL metadata. The compiler does not derive semantic meaning from the IDL information, and fails
15 to fully utilize combined knowledge of IDL and programming language code.

SUMMARY OF THE INVENTION

 The present invention overcomes these problems by unifying a programming language with a definition language. This unification creates a
20 powerful tool for object-oriented programming.

 A programming language is extended to include constructs for a definition language. This simplifies programming by providing an intuitive, easy to use, and systematic way to write programs with definition language semantics. A

programmer places definition language constructs in a programming language file, writing less code and saving time and money. The combined definition language and programming language uses a single set of rules, which simplifies editing and understanding a file.

- 5 In a compiler environment, a compiler for the programming language with embedded definition language information recognizes the definition language constructs. The compiler environment unifies representation of the definition language and programming language, and utilizes definition language information which otherwise would be stripped by a definition language compiler. This
- 10 enables powerful object-oriented programming capabilities within the compiler environment.

- The compiler environment checks for errors in the definition language information and programming language code, and checks for errors caused by inconsistency between definition language information and programming language
- 15 code. Catching errors during development prevents unexplained crashes at run time. In addition, the compiler environment makes the build process more robust, improves error handling, and simplifies debugging during development.

- Processing definition language information along with the programming language code simplifies and speeds up the build process. A programmer
- 20 prepares a single file, and the compiler outputs a single file from which a linker creates an appropriate binary file.

Additional features and advantages of the invention will be made apparent from the following detailed description of an illustrative embodiment that proceeds with reference to the accompanying drawings.

5

BRIEF DESCRIPTION OF THE DRAWINGS

Figure 1 is a block diagram of a software object and a client that interact across interfaces of the software object according to the prior art.

Figures 2a and 2b are a listing for an IDL file according to the prior art.

Figure 3 is a block diagram of a programming language compiler, an IDL
10 compiler, and a linker that create a binary file according to the prior art.

Figure 4 is a block diagram of a computing environment in which the illustrated embodiment can be implemented.

Figure 5 is a block diagram of a software object described in the source code listings of Figures 6a and 6b.

15 Figures 6a and 6b are a source code listing for a file having C++ source code with embedded IDL information describing the software object of Figure 5.

Figure 7 is a block diagram of a C++ compiler that compiles C++ source code with embedded IDL in conjunction with an IDL attribute provider.

Figure 8 is a flow chart showing compilation of a file having C++ source
20 code with embedded IDL information.

Figures 9a-9f are a source code listing representative of the output of the compiler of Figure 7 upon processing the source code listing of Figures 6a and 6b.

Figures 10a and 10b are block diagrams of a C++ compiler and a linker that create a binary file.

DETAILED DESCRIPTION OF AN ILLUSTRATIVE EMBODIMENT

5 The illustrative embodiment of the present invention is directed to unifying an interface definition language ["IDL"] with the C++ programming language. This unification enables powerful interface-based programming capabilities in a C++ compiler environment.

10 The C++ language is extended to include IDL constructs. In the C++ compiler environment, a C++ compiler recognizes the IDL constructs. The C++ compiler environment unifies representation of the embedded IDL and C++ code in a parse tree, and utilizes marshaling, type, and interface information in the IDL which otherwise would be stripped by an IDL compiler. For example, the C++ compiler environment can determine which arguments of an interface are "in,"
15 "out," or "retval" arguments, determine dispatch interface information such as dispids, or determine whether an interface is late or early bound (viz., dispinterface or custom/dual).

 The C++ compiler environment derives semantic meaning from the IDL embedded in C++ code, which enables powerful interface-based programming
20 capabilities. For example, the C++ compiler can automatically generate an implementation for a dispatch interface using an extremely clean and sparse syntax. The C++ compiler can also automatically generate code for a client-side dispatch interface call site, which greatly simplifies programming.

In alternative embodiments, the present invention uses a programming language other than C++ or definition language information expressed in a language other than COM IDL. The programming language can be any language that allows a programmer to express computer processing instructions for an object in a symbolic manner largely independent of machine-specific details. The definition language information can be any type of declarative information that describes an object or an interface exposed by an object.

Although the illustrative embodiment describes COM, alternative embodiments treat other object models. In general, an object model is a set of conventions governing how objects are structured and interact with other objects.

Instead of a C++ compiler environment, alternative embodiments use other programming environments that derive semantic meaning from definition language information embedded in programming language code. A programming environment includes one or more development tools used to develop objects.

For example, in one alternative embodiment, a Java compiler environment facilitates the creation of Java objects (e.g., JavaBeans) that are coded in Java with embedded Java IDL. In another alternative embodiment, a programming environment facilitates the creation of CORBA-compliant objects described in programming language code with embedded CORBA IDL.

In general, the present invention can be applied to any software object for which programming language code includes definition language constructs.

I. Exemplary Computing Environment

Figure 4 illustrates a generalized example of a computing environment 400 in which the illustrative embodiment may be implemented. Computing environment 400 is not intended to suggest any limitation as to scope of use or
5 functionality of the invention, as the present invention may be implemented in diverse general purpose or special purpose computing environments.

With reference to Figure 4, computing environment 400 includes at least one processing unit 410 and memory 420. In Figure 4, this most basic configuration is included within dashed line 430. The processing unit 410
10 executes computer executable instructions and may be a real or a virtual processor. In a multi-processing system, multiple processing units execute computer executable instructions to increase processing power. The memory 420 may be volatile memory (e.g., registers, cache, RAM), non-volatile memory (e.g., ROM, EEPROM, flash memory, etc.), or some combination of the two.

15 A computing environment may have additional features. For example, the computing environment 400 includes storage 440, one or more input devices 450, one or more output devices 460, and one or more communication connections 470. A bus, controller, network, or other interconnection mechanism (not shown) interconnects the components of the computing environment 400.
20 Typically, operating system software (not shown) provides an operating environment for other software executing in the computing environment 400, and coordinates activities of the components of the computing environment 400.

The storage 440 may be removable or non-removable, and includes magnetic disks, magnetic tapes or cassettes, CD-ROMs, DVDs, or any other medium which can be used to store information and which can be accessed within the computing environment 400.

5 An input device 450 may be a touch input device such as a keyboard, mouse, pen, or trackball, a voice input device, a scanning device, or another device that provides input to the computing environment 400. An output device 460 may be a display, printer, speaker, or another device that provides output from the computing environment 400.

10 A communication connection 470 enables communication over a communication medium to another computing entity. The communication medium conveys information such as computer executable instructions or other data in a modulated data signal. A modulated data signal is a signal that has one or more of its characteristics set or changed so as to encode information in the signal. By
15 way of example, communication media include wired or wireless techniques, techniques implemented with an electrical, optical, RF, infrared, acoustic, or other carrier, and combinations of the above.

The invention can be described in the context of computer readable media. Computer readable media are any available media that can be accessed within a
20 computing environment. With the computing environment 400, for example, computer readable media include the memory 420, the storage 440, the communication media, and combinations of the above.

The invention can be described in the context of computer executable instructions (such as those included in program modules) being executed in a computing environment on a target real or virtual processor. Generally, program modules include routines, programs, libraries, objects, classes, components, data structures, etc. that perform particular tasks or implement particular abstract data types. The functionality of program modules may be combined or split between program modules as desired in various embodiments. Computer executable instructions for program modules may be executed within a local or distributed computing environment.

II. Definition Language Information Embedded in Source Code

According to the illustrated embodiment, IDL constructs are embedded in C++ code. The embedded IDL information provides interface definitions and type information for an object, as well as marshaling information. The embedded IDL information includes information that would conventionally be stripped by an IDL compiler.

In alternative embodiments, instead of COM IDL information, the definition language information is Java IDL information, CORBA IDL information, or another type of declarative information that describes an object or an interface of an object. Moreover, the programming language can be Java, C, Basic, or another programming language that allows a programmer to express computer processing instructions for an object in a symbolic manner largely independent of machine-specific details.

Embedded IDL constructs are identified by brackets “[” and “]” marking the beginning and end of a set of one or more IDL attributes, respectively. The set of IDL attributes along with brackets is termed an IDL attribute tag. Embedded IDL in C++ source code for a function declaration Grade of a dual interface ITest is:

```
5      [dual, uuid(12341234-1234-1234-1234-123412341234)] __interface
      ITest {
          [id(1), propget] HRESULT Grade([in] int x, [out, retval] BSTR* pb);
      };
```

10 The IDL attribute “dual” annotates the interface ITest. In general, an IDL attribute tag can be attached to any C++ programming construct, including an entire block or program. Table 1 lists some representative forms.

<i>C++ Construct</i>	<i>Sample</i>
Program (anonymous)	[IDLAttr];
Block	[IDLAttr] { <statements> }]
Class	[IDLAttr] struct A
Interface	[IDLAttr] __interface ITest
Function	[IDLAttr] void Grade
Function argument	void Grade ([IDLAttr] int n)
Variable	[IDLAttr] int n

Table 1

An example of IDL attribute-tag syntax is:

```
15      [IDLAttr1, IDLAttr2 = val1, IDLAttr3(val2), IDLAttr4(prop1 = val3, prop2 =
      val4), IDLAttr5(val5,val6,prop3 = val7,prop4 = {val8,val9}) = {val10,val11}].
```

Within the paired brackets denoting an IDL attribute tag, a first IDL attribute has only a name. A second IDL attribute is set equal to a specific value. A third IDL attribute has a value passed to it as an argument, which is equivalent to the second form. A fourth IDL attribute contains a list of properties set to

20

specific values. The last IDL attribute illustrates a parameter list containing values, properties set to individual values, and a property set to a list of values. Finally, a parameterized IDL attribute itself is set to a list of values. Commas separate multiple attributes in a single tag.

5 Figure 5 illustrates an example of an object CTest 500 that exposes the interfaces ITest 510 (defined in the code portion 630 of Figure 6a) and ITest2 520 (defined in the code portion 640 of Figure 6b). The interface ITest 510 is a dispatch interface that includes methods for invoking the methods Grade and Score. The interface ITest2 includes the methods Display and Hours, as well as
10 methods for retrieving and setting the property StudentID. The object 500 and interfaces 510 and 520 are defined and implemented with a file 600 having C++ source code with embedded IDL that is shown in Figures 6a and 6b. The embedded IDL in the file 600 roughly corresponds to the IDL presented in Figures 2a and 2b. Alternatively, different IDL tags or different programming language
15 constructs can convey equivalent definition information.

The file 600 includes both definitions and implementations of the interfaces ITest 510 and ITest2 520 of the object CTest 500. To facilitate this integration, the usage of embedded IDL is made to be as close as possible to the usage of conventional IDL. For IDL constructs that do not map cleanly to C++, however,
20 deviation from conventional IDL usage is required. Except where otherwise indicated below, embedded IDL constructs are used as in conventional IDL.

The file 600 includes numerous embedded IDL constructs. For example, the code portion 610 includes a construct with the embedded IDL attribute

"project." Rather than use the conventional IDL library block construct in C + +, a global anonymous "project" attribute specifies type library characteristics, including type (e.g., dll or executable), name, helpfile, and helpcontext.

The code portion 620 includes two examples of the IDL attribute "export."

- 5 User-defined data types (struct, enum, union, etc.) that should be visible to embedded IDL are annotated with the "export" attribute. In conventional IDL, these user-defined types were separately specified in an IDL file.

The code portion 630 includes IDL attributes for an interface, for interface member functions, and for arguments. The new keyword "__interface" marks an
10 interface. The "__interface" keyword semantically is a struct which contains pure virtual function members (virtual function members for which no implementation is specified that must be overridden in a derived class in order to create an object). When an interface is annotated with one of the IDL attributes "object," "dispinterface," or "dual," the interface becomes a COM interface (i.e., one that is
15 compliant with Microsoft Corporation's Component Object Model ["COM"]). Accordingly, methods are standard call (__stdcall) by default (and any virtual overrides will also be __stdcall by default). The interface is visible to embedded IDL within it. If a COM interface is specified with no base class, then "object" causes inheritance from IUnknown, whereas "dual" and "dispinterface" cause
20 inheritance from IDispatch. For more information about COM and the IDispatch interface, see Kraig Brockshmidt, Inside OLE, second edition, Chapters 14 and 15, Microsoft Press (1995). In alternative embodiments, the objects described with

embedded definition language information conform to another object model or expose other types of interfaces.

The code portion 630 includes IDL attributes that annotate the interface ITest. ITest is defined as a dispinterface with a helpstring "interface ITest." The
 5 code portion 630 also includes embedded IDL attributes for member functions of ITest. Since ITest is a dispinterface, member functions are given dispids with the "id" IDL attribute. Arguments for member functions are annotated with directional IDL attributes "in," "out," and "retval."

The code portion 640 includes IDL attributes that annotate the interface
 10 ITest2. The "object" IDL attribute indicates ITest2 is a standard COM interface. ITest2 has an associated character helpstring "interface ITest2." The "library_block" IDL attribute indicates that ITest2 should be defined within the library block rather than outside the library block. In contrast, as a dispinterface, ITest is always defined in the library block. The "propput" and "propget" IDL
 15 attributes mark interface members as responsible for setting and retrieving a property value StudentID, respectively. Again, directional IDL attributes "in," "out," and "retval" annotate arguments.

The code portion 650 includes IDL attributes that annotate the object CTest. The IDL attribute "coclass" annotates the struct CTest as a COM class.
 20 Compared to conventional usage, the illustrated embodiment uses a "coclass" IDL attribute to annotate a C++ struct rather than have a coclass definition in an IDL file. The code portion 650 also includes a "helpstring" IDL attribute and a "uuid" IDL attribute. The "uuid" IDL attribute provides a 128-bit unique identifier for the

coclass. The "uuid" IDL attribute can mark other declarations, including interfaces and modules. The "progid" IDL attribute specifies a version dependent alias CTest.17 for the uuid. In the registry, the alias CTest.17 will point to the uuid for CTest. If a uuid, progid, vi_progid, or version is needed for a declaration but not

5 specified, one will be automatically generated in a reproducible fashion.

Finally, the code portion 660 includes implementation code in C + + for the coclass CTest that exposes ITest and ITest2.

As noted above, the illustrated embodiment introduces or changes the usage of several IDL attributes. These special case IDL attributes typically deviate

10 from conventional IDL usage because of C + + incompatibility. Table 2 explains the usage and significance of these special case IDL attributes.

IDL attribute	Usage and significance
coclass	A class implementing interfaces is annotated with the "coclass" IDL attribute and is thus visible as a conventional IDL coclass. In conventional IDL usage, a separate IDL file includes the coclass definition.
cpp_quote	The "cpp_quote" IDL attribute maps to its IDL counterpart in a generated IDL file. In conventional IDL, cpp_quote is a preprocessor directive.
emitidl	By default, the compiler emits IDL metadata to output files. If "emitidl" is set to false, the compiler does not emit IDL information.
export	A user-defined data type annotated with the "export" IDL attribute is visible as a conventional IDL type definition. In conventional IDL, a separate IDL file includes these user-defined type definitions.
idl_module	The "idl_module" IDL attribute is used to import a module such as a dynamic link library in a manner analogous to the conventional IDL attribute "module."
idl_quote	An IDL attribute marked with the "idl_quote" IDL attribute passes through the compiler as IDL metadata. This allows the compiler to handle new IDL attributes for which the compiler lacks full support.
import	The "import" IDL attribute maps to its IDL counterpart in a generated IDL file. In conventional IDL, import is a preprocessor directive.

importidl	The "importidl" IDL attribute automates the determination of whether to import an IDL file to the top of a file or to the beginning of a library block, or to place parts of the IDL file in both.
importlib	The "importlib" IDL attribute maps to its IDL counterpart in a generated IDL file. In conventional IDL, importlib is a preprocessor directive.
include	The "include" IDL attribute maps to its IDL counterpart in a generated IDL file. In conventional IDL, #include is a preprocessor directive.
includelib	The "includelib" IDL attribute places an imported IDL at the beginning of a library block rather than at the top of a file. This facilitates importation of legacy IDL files.
library_block	The "library_block" IDL attribute annotates an interface, struct, enum, etc. to define the annotated feature inside the library block. Otherwise, the annotated feature is defined outside the library block with other marshaling-related information, which may cause an error in some situations.
pragma	The "pragma" IDL attribute maps to its IDL counterpart in a generated IDL file. In conventional IDL, pragma is a preprocessor directive.
progid	The "progid" IDL attribute allows specification of an alias for a CLSID. In conventional usage, a registry script specifies progid.
project	The "project" IDL attribute specifies type library characteristics. In conventional IDL, the block annotated with the attribute "library" includes information used to generate a type library.
satype	The "satype" IDL attribute specifies a type for a SAFEARRAY. For example, satype(BSTR) signifies a SAFEARRAY of BSTRs.
threading	The "threading" IDL attribute specifies a threading model to be used for a file. In conventional usage, a registry script specifies the threading model.
vi_progid	The "vi_progid" IDL attribute allows specification of a version independent progid. This vi_progid remains the same as the progid changes between versions. In conventional usage, a registry script was used to update vi_progid to point to the latest progid.

Table 2

Aside from the IDL attributes presented in Table 2, the illustrated embodiment supports numerous attributes in a manner analogous to conventional usage. Table 3 presents these IDL and ODL attributes. For more information on the usage of these attributes, see Al Major, COM IDL & Interface Design, Wrox Press (1999).

IDL Attributes				ODL Attributes
appobject	helpfile	nonextensible	restricted	aggregatable
bindable	helpstring	object	retval	defaultcollelem
call_as	hidden	odl	size_is	defaultvtable
case	id	oleautomation	string	helpstringcontext
control	ignore	optional	source	helpstringdll
custom	iid_is	out	switch_is	nonbrowsable
default	immediatebind	propget	switch_type	noncreatable
defaultbind	in	propput	transmit_as	range
defaultvalue	last_is	propputref	unique	replaceable
dispinterface	lcid	pointer_default	usesgetlasterror	uidefault
displaybind	length_is	ptr	uuid	
dual	licensed	public	v1_enum	
entry	local	readonly	version	
first_is	max_is	ref	vararg	
helpcontext	ms_union	requestedit	wire_marshal	

Table 3

A programmer can create a file having C++ source code with embedded IDL with the assistance of programming tools. Alternatively, a programmer can create such a file line by line.

5

III. Programming Environment

Figure 7 shows a compiler environment 700 for recognizing embedded IDL constructs in C++ source code, creating a unified representation of the embedded IDL and C++ source code in a parse tree, and deriving semantic meaning from the embedded IDL. The compiler environment 700 also provides error detection for C++ source code with embedded IDL information. In alternative embodiments, another type of programming environment is used to process programming language code with embedded definition language information. For example, in one alternative embodiment, a Java compiler

10

environment processes Java source code with embedded Java IDL to produce virtual machine instructions (e.g., bytecodes) for a Java object.

The compiler environment 700 includes a C++ compiler 720 that accepts a file 710 having C++ source code with embedded IDL. The compiler

5 environment 700 processes the file 710 in conjunction with one or more IDL attribute providers 770, and outputs an output code file 790. Although Figure 7 depicts a single IDL attribute provider 770, the compiler 720 can work with multiple IDL attribute providers (e.g., different providers for different IDL constructs). Alternatively, the functionality of the provider 770 can be merged
10 with the compiler 720.

A front end module 722 reads and performs lexical analysis upon the file 710. Basically, the front end module 722 reads and translates a sequence of characters in the file 710 into syntactic elements, or "tokens," indicating constants, identifiers, operator symbols, keywords, punctuation, etc.

15 A converter module 724 parses the tokens into an intermediate representation. For tokens from C++ source code, the converter module 724 checks syntax and groups tokens into expressions or other syntactic structures, which in turn coalesce into statement trees. Conceptually, these trees form a parse tree 732. As appropriate, the converter module 724 places entries into a
20 symbol table 730 that lists symbol names and type information used in the file 710 along with related characteristics. A symbol table entry for a particular symbol can have a list of IDL attributes associated with it.

If processing the listing of Figures 6a and 6b, for example, the compiler 720 recognizes brackets as delimiting an "export" IDL attribute to be processed as a unit. The following keyword enum denotes a user-defined type element. Within that element is definition for enum E, and so forth. In the symbol table 730, an entry for enum E has the "export" IDL attribute associated with it.

A state 734 tracks progress of the compiler 720 in processing the file 710 and forming the parse tree 732. For example, different state values indicate that the compiler 720 has encountered an IDL attribute, is at the start of a class definition or a function, has just declared a class member, or has completed an expression. As the compiler 720 progresses, it continually updates the state 734. The compiler 720 may partially or fully expose the state 734 to an outside entity such as the provider 770, which can then provide input to the compiler 720.

Based upon the symbol table 730 and the parse tree 732, a back end module 726 translates the intermediate representation of file 710 into output code. The back end module 726 converts the intermediate representation into instructions executable in a target processor, into memory allocations for variables, and so on. In Figure 7, the output code is executable in a real processor, but in alternative embodiments the output code is executable in a virtual processor.

The front-end module 722 and the back-end module 726 can perform additional functions, such as code optimization, and can perform the described operations as a single phase or multiples phases. Except as otherwise indicated,

the modules of the compiler 720 are conventional in nature, and can be substituted with modules performing equivalent functions.

In Figure 7, the provider 770 indicates how to integrate tokens for IDL constructs into the intermediate representation, for example, adding IDL attributes to a symbol table entry for a particular symbol or manipulating the parse tree 732. Thus, embedded IDL information is associated with logically proximate programming language information in the symbol table 730 and the parse tree 732.

In Figure 7, the provider 770 is external to the compiler 720, and communicates with the compiler 720 across the interfaces 750 and 780. Figure 7 depicts a simplified interface configuration of the interface 750 exposed by the compiler 720 and the interface 780 exposed by the provider 770. Alternative interface configurations are possible.

The provider 770 includes several modules. An input module 772 receives a particular IDL attribute from a tag and parses it for parameters, values, properties, or other specifications. The interfaces 750 and 780 define how this information is passed between the compiler 720 and the provider 770.

An operation module 774 determines what must be done to implement the IDL attribute, and identifies locations where code is to be injected, or where other operations are to be carried out. "Injected code" typically includes added statements, metadata, or other elements at one or more locations, but this term also includes changing, deleting, or otherwise modifying existing source code. Injected code can be stored in the provider 770 as one or more templates 776, or

in some other form. In addition, parse tree transformations may take place, such as altering the list of base classes or renaming identifiers.

An output module 778 communicates back to the compiler 720 to effect changes based upon the IDL attributes. In Figure 7, the output module 778
5 directly manipulates internal compiler structures such as the symbol table 730 and the parse tree 732, creating symbols, adding to the parse-tree, etc. Alternatively, the output module 778 writes injected code to an external file (not shown) or send code directly to the compiler 720 as a stream of bytes (not shown) that the compiler 720 processes. Having an IDL attribute provider instruct the compiler
10 (e.g., at converter module 724) to perform the operations gives a measure of security -- the compiler 720 can reject or modify any request that would compromise proper functioning.

As the front end module 722 encounters IDL attribute tags in the file 710, the compiler 720 changes the state 734 appropriately and saves the IDL attribute
15 tags in a list. This list also identifies the location of the provider 770 or any other needed attribute provider, as necessary acquiring location information from a utility such as a registry.

The compiler 720 communicates the state 734 to the provider 770. When the provider 770 detects a point at which it desires to perform an operation, it
20 signals the compiler 720 and effects changes in one of the ways mentioned above. Thus, based upon the semantics of the embedded IDL, the provider 770 affects the states and structures of the compiler 720.

At various points during the processing of the file 710, an error checker module 740 checks for errors in the C++ source code with embedded IDL. In conjunction with the front end module 722, the error checker module 740 detects errors in lexical structure of C++ source code tokens and embedded IDL tokens.

- 5 With converter module 724, error checker 740 detects any syntactical errors in the organization of C++ source code tokens and embedded IDL tokens. The error checker module 740 can also flag certain semantic errors in the embedded IDL in the C++ source code with embedded IDL.

Figure 8 shows a technique 800 for processing IDL embedded in C++ source code in a compiler environment such as that shown in Figure 7.

Alternatively, technique 800 can be performed by a different configuration of modules.

After a compiler reads in (act 810) the file 805, the elements of the file 805 are processed. The compiler gets (act 820) a syntactic element of the file 805 and sets (act 822) an appropriate state for that element. The compiler determines (act 824) whether that element is for a conventional C++ construct or for an IDL construct.

If the current element is for a C++ construct, the compiler converts (act 826) that element into an intermediate language. As appropriate, the compiler handles (act 828) the element, for example, by placing a node in the parse tree or adding an entry to the symbol table. If the compiler determines (act 830) that the file 805 includes more elements, the compiler proceeds with the next element.

If the current element is for an IDL construct, the compiler gets attribute information for the IDL construct. The compiler calls (act 840) an IDL attribute provider, transmitting any parameters or other data accompanying the attribute in the construct. The IDL attribute provider parses (act 850) the passed IDL attribute information.

The IDL attribute provider executes concurrently with the compiler, and more than one IDL attribute provider can be active and executing concurrently during compilation. The IDL attribute provider is loaded upon encountering the first IDL attribute, and stays loaded for the duration of the compile operation. In this way, the IDL attribute provider acts as a domain-specific compiler plug-in that is called to parse constructs that are "registered" as part of its domain.

While executing concurrently with the compiler, the IDL attribute provider detects (act 852) the occurrence of designated events within the compiler, for example, events relate to the state of compilation (in Figure 8, represented by a dashed arrow from act 822 to act 852). The compiler exposes a compilation state to the IDL attribute provider. Examining the state, the IDL attribute provider determines whether to do nothing or to perform (act 854) an operation. Thus, the IDL attribute provider can wait until the compiler reaches a certain state, and then perform an operation when that state is reached, for example, requesting the compiler to modify the parse tree. The IDL attribute provider then waits for another event.

The IDL attribute provider can perform different operations for different events that might occur within the compiler, and for different parameters

transmitted with an IDL attribute. Among these operations are injection of statements or other program elements, possibly employing templates, and modifying or deleting code. Other operations include adding new classes, methods and variables, or modifying existing ones. Modification can include

5 renaming and extending an object or construct. In Figure 8, dashed arrows from act 854 to acts 826 and 828 represent the passing of code, state information, instructions, or other data to the compiler as described above.

Injected code is typically located remotely from where the IDL attribute appears in the C++ source code. Code can be injected at multiple locations as

10 well. To clarify the significance of the injected code, comments around the injected code can identify the IDL attribute for which it was injected.

The scope of an IDL attribute is not bound to the scope of its associated C++ construct (variable, class, etc., see Table 1). Rather, the scope of an IDL attribute can extend beyond the point of its use. In most cases, however, an IDL

15 attribute affects semantics in the context of its associated C++ construct. In Figure 6a, for example, the "export" IDL attributes are coextensive with the scope of their respective type definitions, and the "dispinterface" IDL attribute operates over the ITest interface declaration.

When the file 805 has been completely processed, the compiler translates

20 (act 870) the intermediate representation into output code 890 such as machine-language instructions, memory allocations, and other data required for a target processor. When the compiler finishes the compile operation, the IDL attribute provider exits (act 860).

IDL and programming languages evolve independently, and IDL typically changes more rapidly. Accordingly, using the "idl_quote" IDL attribute, the compiler can pass through new IDL attributes as IDL metadata for subsequent processing by an IDL compiler that understands the IDL attributes.

5 Although Figure 8 depicts acts in a particular order, per conventional compiler techniques, many of these acts can be rearranged or performed concurrently. For example, the acts of reading the file, getting elements, determining significance, and translating to output code can be overlapped to some extent.

10 Figures 9a through 9f show a source code listing 900 representative of the output of the compiler environment of Figure 7 upon processing the source code listing of Figures 6a and 6b. Because the listing 900 is in source code form, it is merely representative of the compiler output, which can be and typically is in a computer executable form.

15 Figures 9a through 9f show a representation of a compiler-generated implementation of an IDispatch interface for the interface ITest, which is described in the code portions 630 and 660 of Figures 6a and 6b. For the sake of presentation, the original C++ and bracketed IDL attributes are presented in boldface. The code generated based upon the semantics of the embedded IDL in
20 the C++ source code is presented in lighter type.

Using conventional techniques, implementing dispatch interfaces in an object is notoriously difficult, as is implementing the client-side code for packing and unpacking variants. According to the illustrated embodiment, the compiler

automatically generates IDispatch interface implementations as well as corresponding client-side code.

The IDispatch interface implementation of Figures 9a through 9f includes the method Invoke for invoking Grade or Score based upon a dispid, packed arguments, and other parameters. The implementation also includes the method
5 GetIdsOfNames to associate the names Grade and Score with dispids, and the methods TypeInfoHelper, GetTypeInfoCount, and GetTypeInfo to provide type information to a client. For more information about IDispatch, see Kraig Brockshmidt, Inside OLE, second edition, Chapters 14 and 15, Microsoft Press
10 (1995).

The compiler environment of Figure 7 can also output client-side code for operating the IDispatch implementation of Figures 9a through 9f. Alternatively, the compiler environment generates other output code based upon combined IDL and programming language semantics.

IV. Creation of Binary Files

Figures 10a and 10b show the creation of a binary file 1090 from a file 1005 having C++ source code with embedded IDL. Figure 10a shows a high-level view of the process 1000 of creating a binary file 1090, while Figure 10b
20 shows one particular implementation 1002. In alternative embodiments, other types of files are created from a programming language file with embedded definition language information. For example, in one alternative embodiment, a

Java compiler environment facilitates creation of a Java file such as a Java Archive ["JAR"] file from a Java source code with embedded Java IDL.

With reference to Figure 10a, a programmer has one or more files 1005. A C++ compiler 1010 converts the one or more files 1005 into one or more output code files 1020 that include fragments of IDL information. Next, a linker 1030 generates a binary file 1090, such as an executable or dynamic link library, from the one or more output code files 1020 with fragments of IDL information. Compared to multi-phase, multi-file techniques, the process 1000 simplifies creation of binary files for the programmer by combining IDL and programming language code into a single file and by automating the process 1000 after the programmer initiates the process 1000.

In one embodiment, based upon the files 1005 the compiler 1010 generates a type library and other output files conventionally generated by an IDL compiler. A drawback to this embodiment, as mentioned above, is that IDL technology historically has evolved apart from compiler technology. Different groups might develop the compiler and the IDL, so the compiler might not be up to date. Still another group might develop the linker.

To let the compiler exploit semantic knowledge of IDL while preserving the independence of the IDL, after deriving semantic meaning from embedded IDL, the compiler can output the embedded IDL as IDL metadata, which is then supplied to an IDL compiler. Figure 10b shows such an embodiment.

In Figure 10b, the linker 1030 takes the fragments of IDL information from the output code files 1020. The linker passes the IDL fragments to the compiler

1010, which generates a collective IDL file 1040. The compiler 1010 embeds debugging information in the IDL file 1040, so subsequent error messages from an IDL compiler can be associated with appropriate embedded IDL in C++ source code.

5 The linker 1030 then calls an IDL compiler 1042, which generates a type library file 1044. The linker 1030 next generates an intermediate resource file 1050 such as an RC file and embeds the type library file 1044 in the intermediate resource file 1050. The linker 1030 calls a resource tool 1052 such as a resource compiler, which generates a resource file 1060 such as a RES file. Next, the
10 linker 1030 calls a resource file combiner 1064 that combines the resource file 1060 and any related resource files (e.g., resource file 1062) into a combined resource file 1070. Taking the combined resource file 1070, the linker 1050 produces a binary file 1090.

 Separating these activities as described in Figure 10b facilitates separate
15 development and control of the various components. On the other hand, many of the activities of Figure 10b could be consolidated without departing from the high level view of Figure 10a. For example, the compiler 1010 could pass a representation of embedded IDL directly to the IDL compiler 1042 instead of placing IDL fragments in the output code files 1020 and subsequently generating
20 the IDL file 1040.

 Having described and illustrated the principles of our invention with reference to an illustrative embodiment, it will be recognized that the illustrative

embodiment can be modified in arrangement and detail without departing from such principles. It should be understood that the programs, processes, or methods described herein are not related or limited to any particular type of computing environment, unless indicated otherwise. Various types of general purpose or specialized computing environments may be used with or perform operations in accordance with the teachings described herein. Elements of the illustrative embodiment shown in software may be implemented in hardware and vice versa.

In view of the many possible embodiments to which the principles of our invention may be applied, we claim as our invention all such embodiments as may come within the scope and spirit of the following claims and equivalents thereto.

CLAIMS

We claim:

1. A computer readable medium having stored thereon a computer executable compiler system that performs semantic analysis of definition language information embedded in programming language code in a file, the compiler system comprising:

a front end module that separates a file into plural tokens, the file including programming language code having embedded therein definition language information;

a converter module that converts the plural tokens into an intermediate representation;

a back end module that produces output code from the intermediate representation based at least in part upon semantics of the embedded definition language information.

2. The compiler system of claim 1 wherein the intermediate representation includes a symbol table and a parse tree that unifies representation of the programming language code and the embedded definition language information.

3. The compiler system of claim 2 wherein the symbol table includes plural entries for symbol names for the programming language code, and wherein

at least one of the plural entries has an associated list of definition language attributes.

4. The compiler system of claim 1 further comprising a definition language attribute provider that modifies the intermediate representation based upon the semantics of the embedded definition language information.

5. The compiler system of claim 1 further comprising an error checker module that checks for lexical, syntactic, and semantic errors in the file.

6. In a computer system, a computer executable compiler system that creates a unified programming language and interface definition language parse tree from a file comprising a mix of programming language constructs and interface definition language constructs, the compiler system comprising:

a front end module that separates a file into plural tokens, the file comprising a mix of programming language constructs and interface definition language constructs;

a converter module that converts the plural tokens into an intermediate representation comprising a symbol table and a parse tree, wherein the symbol table includes plural entries for symbol names for the programming language constructs, at least one of the plural entries having an associated list of interface definition language attributes, and wherein the parse tree unifies representation of

the programming language constructs and the interface definition language constructs.

7. The compiler system of claim 6 wherein the front end module recognizes a delimiting character that distinguishes interface definition language tokens from programming language tokens.

8. The compiler system of claim 6 further comprising an error checker module that performs lexical and syntactic checks on the file.

9. A computer readable medium having stored thereon a data structure representing a unified interface definition language and programming language parse tree for a file having a combination of programming language code and embedded interface definition language information, the data structure comprising:

a first data field storing data representing a symbol table that has plural entries, each of the plural entries corresponding to a symbol name for programming language code of a file having a combination of programming language code and embedded interface definition language information, at least one of the plural entries having an associated list of interface definition language attributes based upon the embedded interface definition language information; and

a second data field storing data representing a parse tree, wherein the parse tree unifies representation of the programming language code and the embedded interface definition language information.

10. In a computer system, a method of creating a binary file from an input file that includes a mix of programming language constructs and definition language constructs, the method comprising:

providing one or more input files, each input file comprising a mix of programming language constructs and definition language constructs;

upon user initiation at compile time, creating a binary file from the one or more input files, wherein the creation of the binary file comprises:

with a compiler, converting the one or more input files into one or more output code files that include fragments of definition language information; and

with a linker, generating a binary file from the one or more output code files.

11. The method of claim 10 wherein the generating comprises:

extracting the fragments of definition language information from the one or more output code files;

passing the extracted fragments to the compiler;

generating by the compiler an intermediate definition language file;

based upon the intermediate definition language file, generating by a definition language compiler a type library file; and

producing the binary file based upon the one or more output code files and the type library file.

12. The method of claim 11 wherein the producing comprises:

- embedding the type library file into a first intermediate resource file;
- with a resource tool, generating a second intermediate resource file;
- with a resource file combiner, combining the second intermediate resource file with one or more related resource files into a combined resource file; and
- producing the binary file based upon the one or more output code files and the combined resource file.

13. In a computer system, a method of deriving semantic meaning from definition language information embedded in programming language code in a file, the method comprising:

- separating a file into plural tokens, the file including definition language information embedded in programming language code;
- converting the plural tokens into an intermediate representation;
- generating output code from the intermediate representation based at least in part upon semantics of the embedded definition language information.

14. The method of claim 13 wherein the converting comprises:

- building a symbol table having plural entries for symbol names for the programming language code, at least one of the plural entries having an associated list of definition language attributes based upon the embedded definition language information; and

building a parse tree that unifies representation of the programming language code and the embedded definition language information.

15. The method of claim 13 further comprising:

modifying the intermediate representation by a definition language attribute provider based upon the semantics of the embedded definition language information.

16. A computer readable medium having stored thereon instructions for performing a method of creating a unified programming language and definition language parse tree from a file that includes definition language information embedded in programming language code, the method comprising:

separating a file into plural tokens, the file including definition language information embedded in programming language code;

building a symbol table having plural entries for symbol names for the programming language code, at least one of the plural entries having an associated list of definition language attributes based upon the embedded definition language information; and

building a parse tree that unifies representation of the embedded definition language information and the programming language code.

17. The computer readable medium of claim 16 wherein the separating comprises recognizing a delimiting character that distinguishes definition language tokens from programming language tokens.

18. A computer readable medium having stored thereon a computer executable compiler system that checks for errors in a file comprising a mix of definition language information and programming language code, the compiler system comprising:

a front end module that separates a file into plural programming language tokens and plural definition language tokens, the file comprising a mix of definition language information and programming language code, wherein the front end module further checks for lexical errors in the file;

a converter module that converts the plural programming language tokens and plural definition language tokens into an intermediate representation, the converter module further checking for syntax errors;

a back end module that produces output code from the intermediate representation based at least in part upon semantics of the definition language information.

19. The compiler system of claim 18 wherein the converter module further checks for semantic errors between the definition language information and the programming language code.

20. In a computer system having a programming language compiler that generates output code based upon programming language source code, the programming language compiler including a compiler state, a symbol table, and a parse tree, an improvement comprising:

modifying the programming language compiler to recognize constructs of interface definition language information embedded within programming language source code;

modifying the programming language compiler to expose the compiler state to one or more interface definition language attribute providers;

modifying the programming language compiler to allow manipulation of the symbol table and the parse tree by the one or more interface definition language attribute providers based upon the semantics of the embedded interface definition language information.

21. In a computer system, a method of embedding debugging information in a definition language output file to facilitate debugging of an input file, the input file comprising constructs of definition language information embedded in programming language code, the method comprising:

receiving by a programming language compiler an input file, the input file comprising constructs of definition language information embedded in programming language code;

embedding by the programming language compiler debugging information in a definition language output file, the definition language output file for

subsequent processing by a definition language compiler, whereby the embedded debugging information associates errors raised by the definition language compiler with locations of embedded definition language constructs in the input file to facilitate debugging of the input file.

UNIFICATION OF A PROGRAMMING LANGUAGE AND A DEFINITION LANGUAGE**ABSTRACT**

Unification of a programming language with a definition language creates a powerful tool for object-oriented programming. A programming language is extended to include definition language constructs. In a compiler environment, a compiler recognizes the definition language constructs in programming language code. The compiler environment unifies representation of the definition language information and programming language code in a unified parse tree, and derives semantic meaning from definition language information. The compiler environment also checks for lexical, syntax, and semantic errors.

FIGURE 1 (Prior Art)

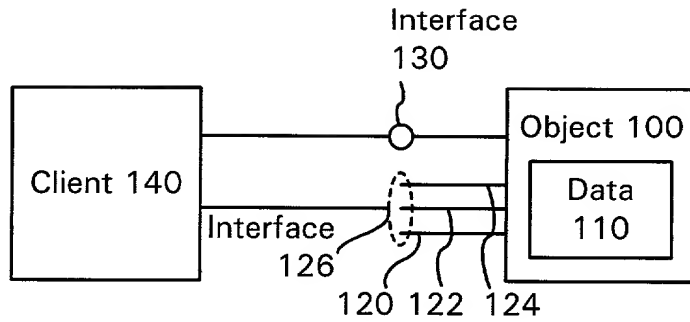


FIGURE 3 (Prior Art)

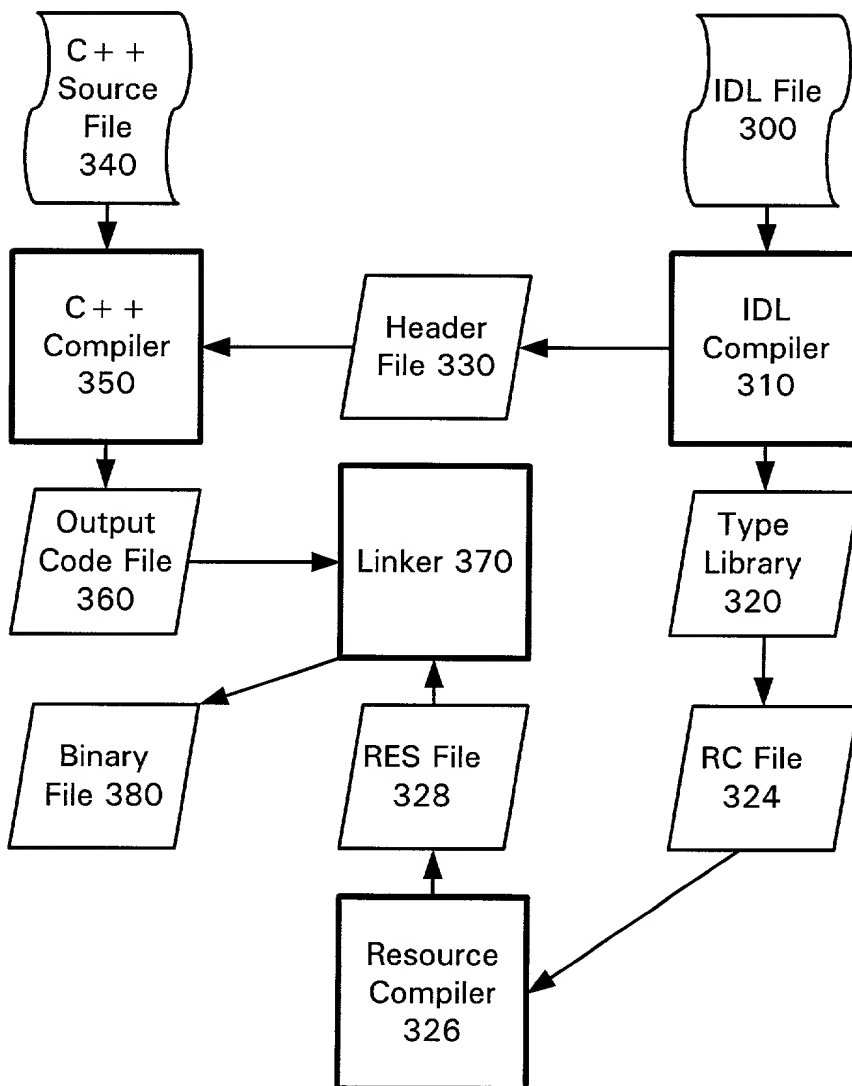


FIGURE 2a (Prior Art)

```
import "docobj.idl";

enum E {
    e = 17,
};

struct S {
    int i, j;
};

[version(1.0), helpfile("test.res"), helpcontext(12), uuid(eed3644c-8488-3ecd-ba97-147db3cddb499) ]
library MyLib
{
    importlib("stdole2.tlb");
    importlib("olepro32.dll");

    [uuid(1AECC9B3-2104-3723-98B8-7CC54722C7DD), helpstring("interface ITest")]
    dispinterface ITest {
        properties:
        methods:
        [id(34)] void Grade([in] enum E , [out,retval] char *);
        [id(18)] HRESULT Score([in] struct S *a, [in] float b, [in] VARIANT c);
    };
}
```

To Figure 2b

FIG. 2a is a diagram of a prior art interface definition language (IDL) file. The diagram shows a code snippet for an IDL file named "docobj.idl". The code defines an enumeration "E" with a value "e = 17", a structure "S" with two integer fields "i" and "j", and a library "MyLib" that imports "stdole2.tlb" and "olepro32.dll". The library "MyLib" defines a dispinterface "ITest" with a uuid "1AECC9B3-2104-3723-98B8-7CC54722C7DD" and a helpstring "interface ITest". The interface "ITest" has two methods: "Grade" which takes an enumeration "E" and returns a character string, and "Score" which takes a structure "S", a float "b", and a variant "c", and returns a HRESULT. The diagram is labeled "FIG. 2a" and "To Figure 2b".

FIGURE 2b (Prior Art)

To Figure 2a

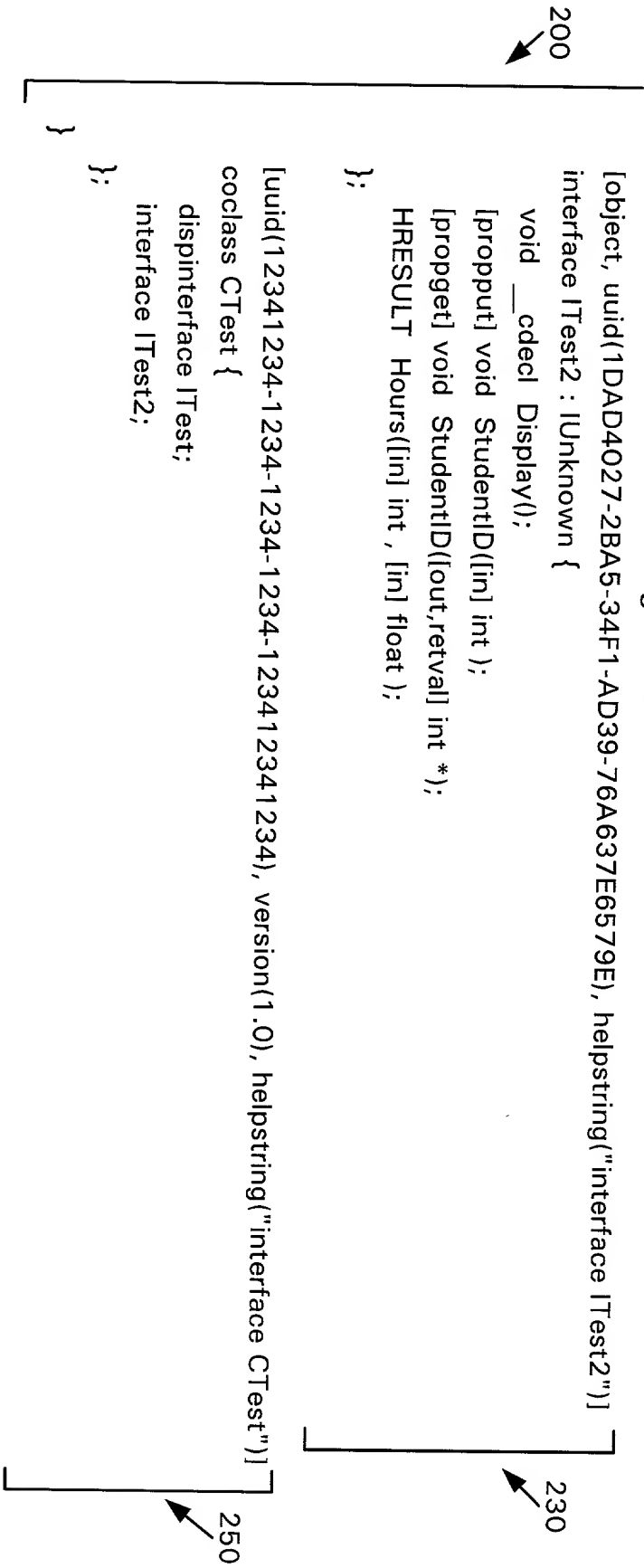


FIGURE 4

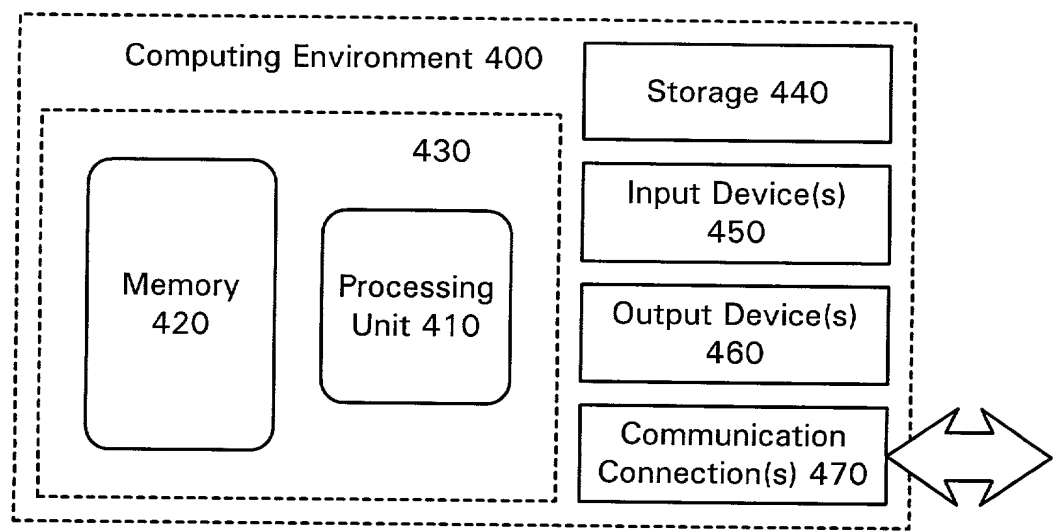
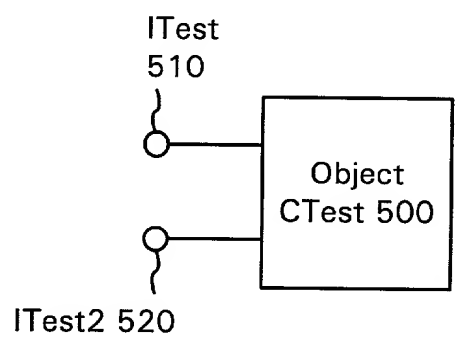


FIGURE 5



600
#define _ATL_ATTRIBUTES 1
#include "atbase.h"
#include "atcom.h"
extern "C" int printf(const char*, ...);

// IDL library block and overall project settings
[project(type=dll, name=MyLib, helpfile="test.res", helpcontext=12)];

[export] enum E {
 e = 17
};
[export] struct S {
 int i, j;
};

[dispinterface, helpstring("interface ITest")] __interface ITest : IDispatch {
 [id(34)] void Grade([in] E, [out, retval] char*);
 [id(18)] HRESULT Score([in]S* a, [in]float b, [in]VARIANT c);
};

[object, library_block, helpstring("interface ITest2")] __interface ITest2 {
 void __cdecl Display(void);
 [propput] void StudentD([in] int);
 [proppet] void StudentID([out,retval] int*);
 HRESULT Hours([in]int, [in]float);
};

To Figure 6b

FIGURE 6a

To Figure 6a

```
600 [ccoclass, progid(CTest.17), helpstring("interface CTest"), uuid(12341234-1234-1234-1234-12341234)]
struct CTest : ITest, ITest2 {
    void Grade(E e, char* pc) {
        printf("CTest::Grade(e=%d)\n", (int) e);
        *pc = 'A';
    }
    HRESULT Score(S* a, float b, VARIANT c) {
        printf("CTest::Score(a=%p,b=%f,c=%d)\n", a, b, c.lVal);
        return S_OK;
    }
    void Display() {
        printf("CTest::Display()\n");
    }
    void put_StudentID(int i) {
        printf("CTest::put_StudentID(i=%d)\n", i);
    }
    void get_StudentID(int* ) {
        printf("CTest::get_StudentID()\n");
    }
    HRESULT Hours(int a, float b) {
        printf("CTest::Hours(a=%d,b=%f)\n", a, b);
        return S_OK;
    }
};
```

FIGURE 6b

FIGURE 7

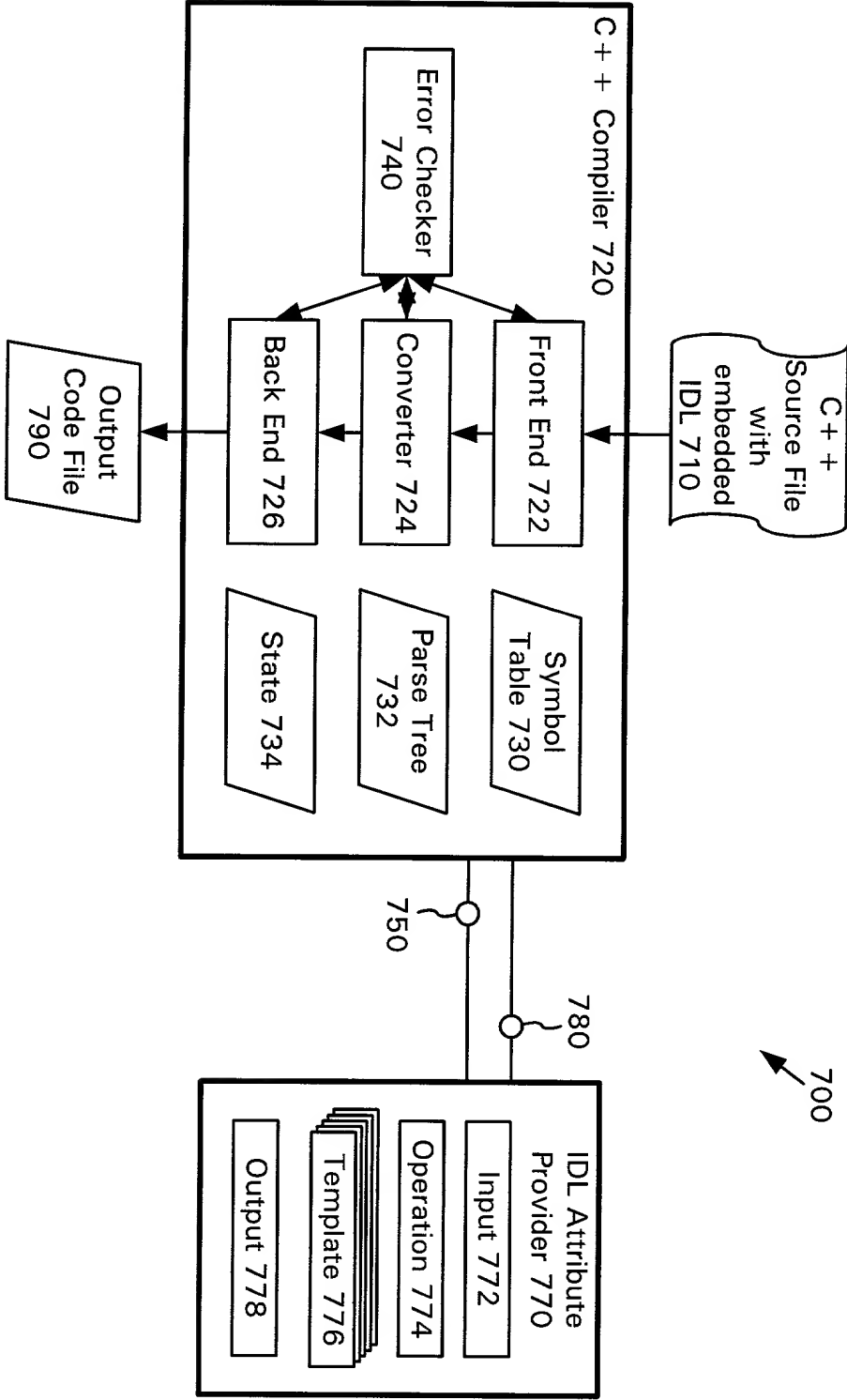
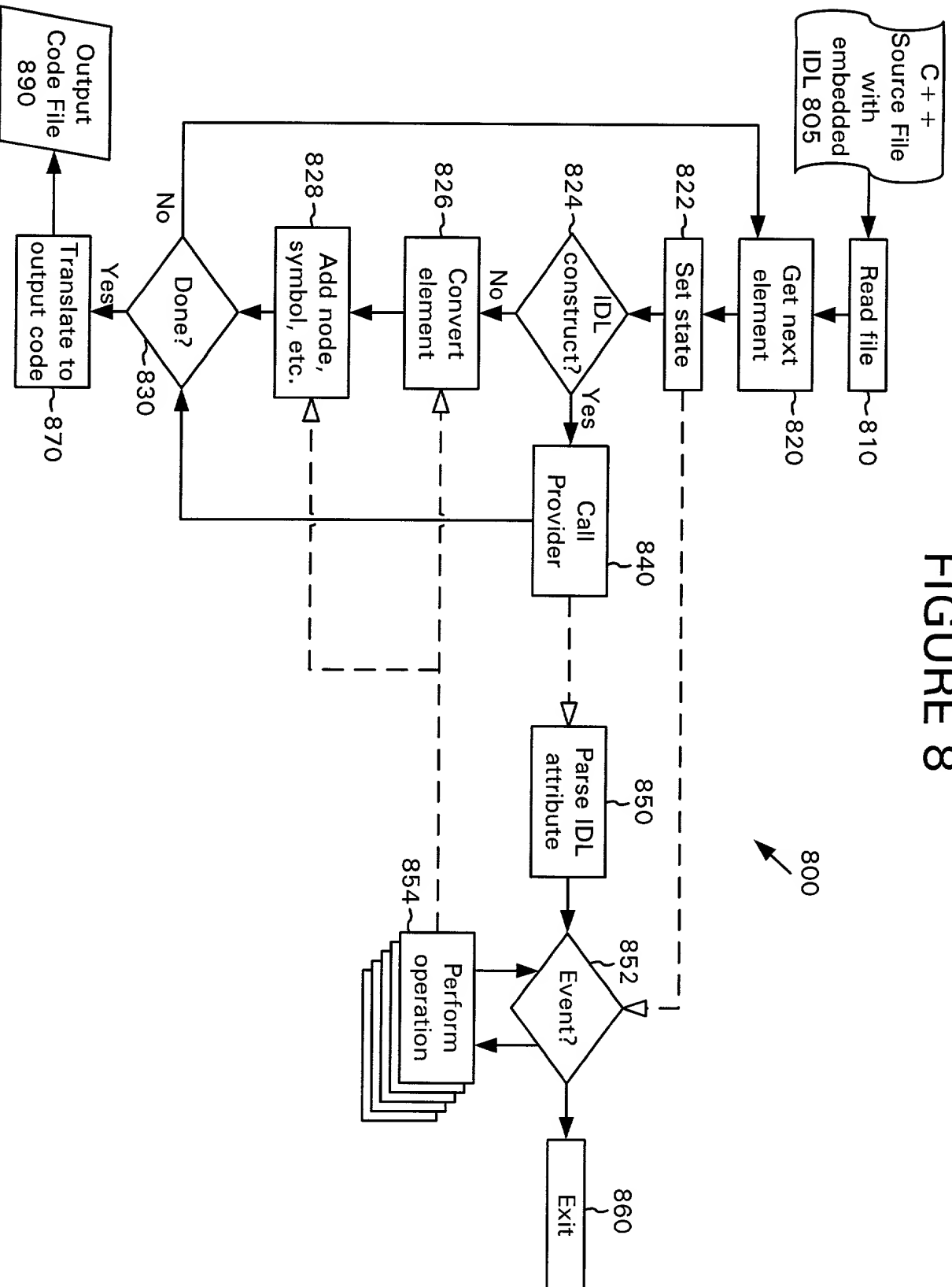


FIGURE 7 is a block diagram of a C++ compiler system 700. The system 700 includes a C++ source file with embedded IDL 710, a C++ compiler 720, and an IDL attribute provider 770. The C++ compiler 720 includes a front end 722, a converter 724, a back end 726, an error checker 740, a symbol table 730, a parse tree 732, and a state 734. The IDL attribute provider 770 includes an input 772, an operation 774, a template 776, and an output 778. The system 700 also includes an output code file 790. The system 700 is connected to a network 750 and a database 780.

FIGURE 8



```

900
[ccclass, progid(CTest.17), helpstring("interface CTest"), uuid(12341234-1234-1234-1234-12341234)]
struct CTest : ITest, ITest2 ,
/* + + + Added Baseclass */ public CComCoClass < CTest, &__uuidof(CTest) >,
/* + + + Added Baseclass */ public CComObjectRootEx < CComSingleThreadModel >,
/* + + + Added Baseclass */ public IProvideClassInfoImpl < &__uuidof(CTest) >
{
    void Grade(E e, char* pc) {
        printf("CTest::Grade(e = %d)\n", e);
        *pc = 'A';
    }
    HRESULT Score(S* a, float b, VARIANT c) {
        printf("CTest::Score(a = %p, b = %f, c = %d)\n", a, b, c.iVal);
        return S_OK;
    }
    void Display() {
        printf("CTest::Display()\n");
    }
    void put_StudentID(int i) {
        printf("CTest::put_StudentID(i = %d)\n", i);
    }
    void get_StudentID(int* ) {
        printf("CTest::get_StudentID()\n");
    }
    HRESULT Hours(int a, float b) {
        printf("CTest::Hours(a = %d, b = %f)\n", a, b);
        return S_OK;
    }
}

```

FIGURE 9a

FIGURE 9a



```
// + + + Start Injected Code
virtual HRESULT STDMETHODCALLTYPE Invoke(
    /* [in] */ DISPID dispIdMember,
    /* [in] */ REFIID riid,
    /* [in] */ LCID lcid,
    /* [in] */ WORD wFlags,
    /* [out][in] */ DISPPARAMS *pDispParams,
    /* [out] */ VARIANT *pVarResult,
    /* [out] */ EXCEPINFO *pExceptionInfo,
    /* [out] */ UINT *puArgErr)
{
    HRESULT hr = S_OK;
    if (pDispParams == 0) {
        return DISP_E_BADVARTYPE;
    }
    if (pVarResult != 0) {
        VariantInit(pVarResult);
    }
    switch (dispIdMember) {
    case 18:
    {
        S * i1 = (S *) V_RECORD(&pDispParams->rgvarg[2]);
        float i2 = V_R4(&pDispParams->rgvarg[1]);
        VARIANT i3 = pDispParams->rgvarg[0];
```

FIGURE 9b


```
virtual HRESULT STDMETHODCALLTYPE GetIdsOfNames(
    /* [in] */ REFIID riid,
    /* [size_is][in] */ LPOLESTR *rgszNames,
    /* [in] */ UINT cNames,
    /* [in] */ LCID lcid,
    /* [size_is][out] */ DISPID *rgDispId)
{
    static LPOLESTR names[] = { L"Grade", L"Score" };
    static DISPID dids[] = { 34, 18 };
    for (unsigned int i = 0; i < cNames; ++i) {
        int fFoundIt = 0;
        for (unsigned int j = 0; j < sizeof(names)/sizeof(LPOLESTR); ++j) {
            if (lstrcmpW(rgszNames[i], names[j]) == 0) {
                fFoundIt = 1;
                rgDispId[i] = dids[j];
            }
        }
        if (fFoundIt == 0) {
            return DISP_E_UNKNOWNNAME;
        }
    }
    return S_OK;
}
```

FIGURE 9d


```

HRESULT TypelInfoHelper(REFIID iidDisp, LCID /*lcid*/, ITypelInfo * * ppTypelInfo)
{
    if (ppTypelInfo == NULL) {
        return E_POINTER;
    }
    *ppTypelInfo = NULL;
    TCHAR szModule1[_MAX_PATH];
    ::GetModuleFileName(_pModule->GetModuleInstance(), szModule1, _MAX_PATH);
    USES_CONVERSION;
    CComPtr<ITypelib> spTypelib;
    HRESULT hr = LoadTypelib(T2OLE(szModule1), &spTypelib);
    if (SUCCEEDED(hr)) {
        CComPtr<ITypelInfo> spTypelInfo;
        hr = spTypelib->GetTypelInfoOfGuid(iidDisp, &spTypelInfo);
        if (SUCCEEDED(hr)) {
            *ppTypelInfo = spTypelInfo.Detach();
        }
    }
    return hr;
}

```

900
↖

FIGURE 9e

```

900
virtual HRESULT STDMETHODCALLTYPE GetTypeInfoCount(unsigned int* pctinfo)
{
    if (pctinfo == NULL) {
        return E_POINTER;
    }
    CComPtr<ITypInfo> spTypeInfo;
    *pctinfo =
        (SUCCEEDED(TypInfoHelper(__uuidof(ITest), 0, &spTypeInfo))) ? 1 : 0;
    return S_OK;
}

virtual HRESULT STDMETHODCALLTYPE GetTypeInfo(unsigned int iTInfo, LCID lcid, ITypInfo** ppTypeInfo)
{
    if (iTInfo != 0) {
        return DISP_E_BADINDEX;
    }
    return TypeInfoHelper(__uuidof(ITest), lcid, ppTypeInfo);
}

// + + + End Injected Code
.....
};

```

FIGURE 9f

FIGURE 10a

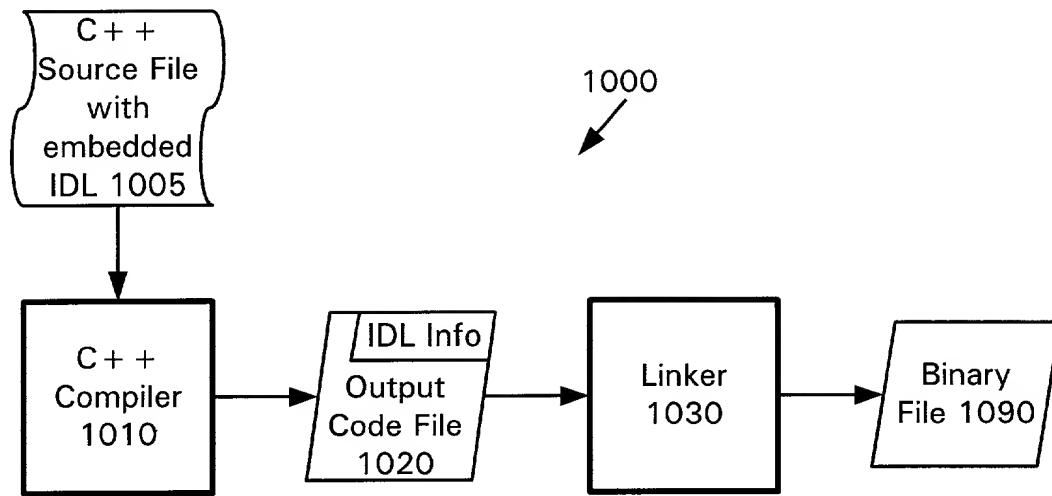
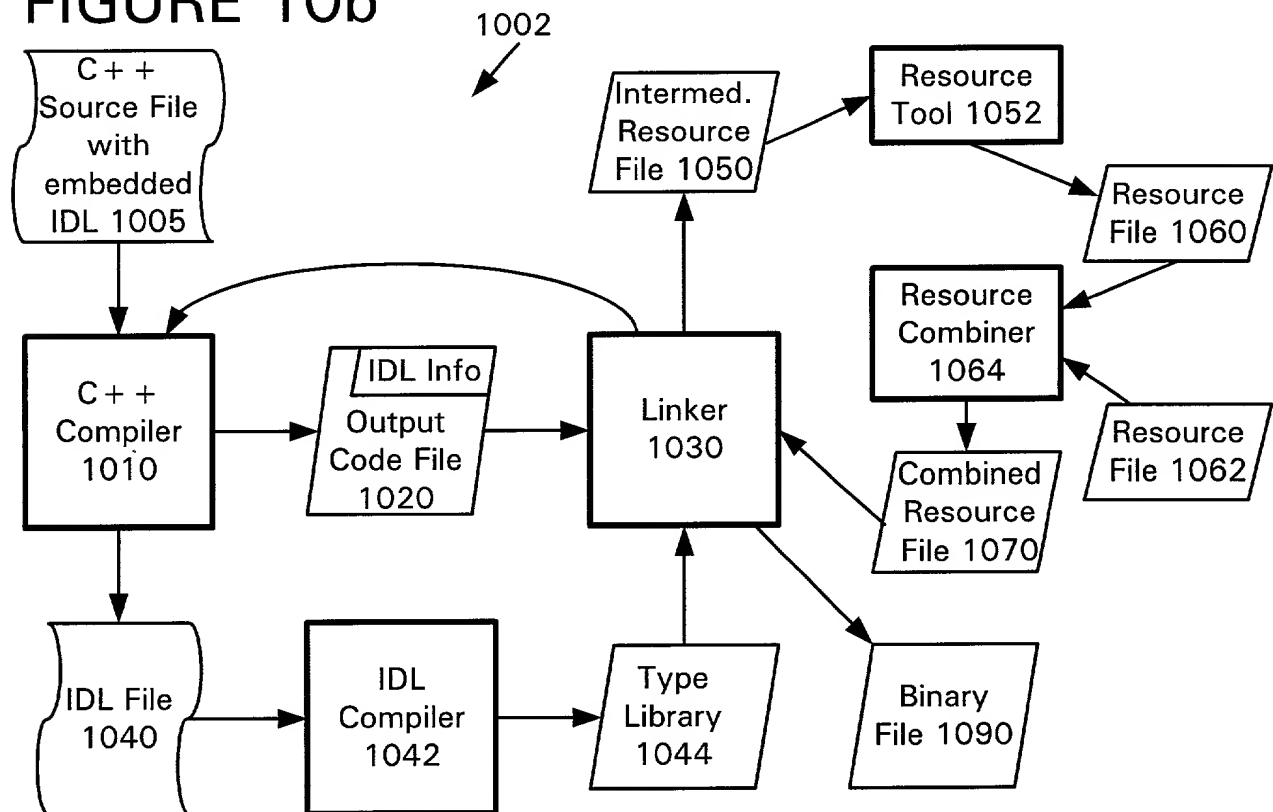


FIGURE 10b



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COMBINED DECLARATION AND POWER OF ATTORNEY FOR PATENT APPLICATION

As a below named inventor, I hereby declare that:

My residence, post office address and citizenship are as stated below next to my name,

I believe I am the original, first and sole inventor (if only one name is listed below) or an original, first and joint inventor (if plural names are listed below) of the subject matter which is claimed and for which a patent is sought on the invention entitled **UNIFICATION OF A PROGRAMMING LANGUAGE AND A DEFINITION LANGUAGE**, the specification of which

☒ is attached hereto.

☐ was filed on as Application No.

☐ was described and claimed in PCT International Application No. _____, filed on _____, and as amended under PCT Article 19 on _____ (if applicable).

☐ and was amended on _____ (if applicable).

☐ with amendments through _____ (if applicable).

I hereby state that I have reviewed and understand the contents of the above-identified specification, including the claims, as amended by any amendment referred to above.

I acknowledge the duty to disclose information which is material to patentability as defined in Title 37, Code of Federal Regulations, § 1.56. If this is a continuation-in-part application filed under the conditions specified in 35 U.S.C. § 120 which discloses and claims subject matter in addition to that disclosed in the prior copending application, I further acknowledge the duty to disclose material information as defined in 37 C.F.R. § 1.56 which occurred between the filing date of the prior application and the national or PCT international filing date of the continuation-in-part application.

I hereby claim foreign priority benefits under Title 35, United States Code, § 119(a)-(d) of any foreign application(s) for patent or inventor's certificate or of any PCT International application(s) designating at least one country other than the United States of America listed below and have also identified below any foreign application(s) for patent or inventor's certificate or any PCT International application(s) designating at least one country other than the United States of America filed by me on the same subject matter having a filing date before that of the application(s) on which priority is claimed:

Prior Foreign Application(s)

Priority
Claimed

_____ (Number)	_____ (Country)	_____ (Day/Month/Year Filed)	<input type="checkbox"/> Yes	<input type="checkbox"/> No
-------------------	--------------------	---------------------------------	---------------------------------	--------------------------------

I hereby claim the benefit under Title 35, United States Code, § 119(e) of any United States provisional application(s) listed below:

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Application Number

Filing Date

I hereby claim the benefit under Title 35, United States Code, § 120 of any United States application(s) or § 365(c) of any PCT International application(s) designating the United States, listed below and, insofar as the subject matter of each of the claims of this application is not disclosed in the prior United States or PCT International application in the manner provided by the first paragraph of Title 35, United States Code, § 112, I acknowledge the duty to disclose material information as defined in Title 37, Code of Federal Regulations, § 1.56(a) which occurred between the filing date of the prior application and the national or PCT International filing date of this application:

(Application No.)

(Filing Date)

(Status: patented,
Pending, abandoned)

The undersigned hereby authorizes the U.S. attorney or agent named herein to accept and follow instructions from _____ as to any action to be taken in the Patent and Trademark Office regarding this application without direct communication between the U.S. attorney or agent and the undersigned. In the event of a change in the persons from whom instructions may be taken, the U.S. attorney or agent named herein will be so notified by the undersigned.

I hereby appoint the following attorney(s) and/or agent(s) to prosecute this application, to file a corresponding international application, and to transact all business in the Patent and Trademark Office connected therewith:

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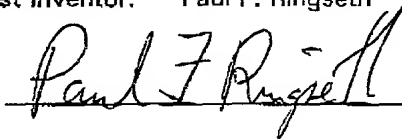
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I hereby declare that all statements made herein of my own knowledge are true and that all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under Section 1001 of Title 18 of the United States Code and that such willful false statements may jeopardize the validity of the application or any patent issued thereon.

Full Name of Sole or First Inventor: Paul F. Ringseth

Inventor's Signature

7/5/00
Date

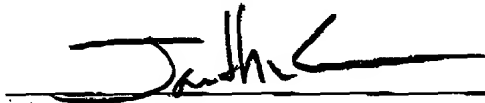
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Inventor's Signature

7/5/00
Date

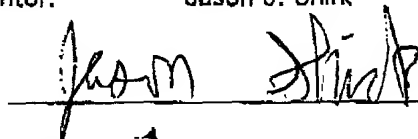
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7-5-00
Date

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